

# ARIPOTER-Degrees: ARgumentation apPrOximaTE Reasoning using In/Out Degrees of Arguments

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**Abstract**—This document describes our approach for approximating the acceptability of arguments under various semantics using their in/out-degrees, and the tool we have implemented in Java. It solves all the problems in the approximate track of ICCMA 2023, namely  $DC\text{-}\{CO|ST|SST|STG|ID\}$  and  $DS\text{-}\{PR|ST|SST|STG\}$ .

**Index Terms**—Abstract argumentation, approximate reasoning, Java tool

## I. BACKGROUND: SEMANTICS AND PROBLEMS

An abstract argumentation framework (AF) [1] is a directed graph  $\mathcal{F} = \langle \mathcal{A}, \mathcal{R} \rangle$  s.t.  $\mathcal{A}$  is the set of *arguments* and  $\mathcal{R} \subseteq \mathcal{A} \times \mathcal{A}$  is the *attack relation* over the arguments. For two arguments  $a, b \in \mathcal{A}$ , the notation  $(a, b) \in \mathcal{R}$  means that  $a$  attacks  $b$ . For a set of arguments  $S \subseteq \mathcal{A}$ , we note  $S^+ = \{a \in \mathcal{A} \mid \exists b \in S, (b, a) \in \mathcal{R}\}$ .

In this work, we focus on extension-based semantics whose aim is to return sets of arguments that are compatible with each other according to different criteria. Mathematically, these semantics are functions  $\sigma$  s.t.  $\sigma(\mathcal{F}) \subseteq 2^{\mathcal{A}}$ . The extension-based semantics that we consider are based on conflict-freeness and admissibility:  $S \subseteq \mathcal{A}$  is *conflict-free* iff  $\forall a, b \in S, (a, b) \notin \mathcal{R}$ , and *admissible* iff it is conflict-free and it defends all its elements, meaning that  $\forall a \in S, \forall b \in \mathcal{A}$  s.t.  $(b, a) \in \mathcal{R}, \exists c \in S$  s.t.  $(c, b) \in \mathcal{R}$ . The semantics are:

- Complete (CO):  $S$  is an extension iff it is an admissible set which does not defend any argument outside of  $S$ ,
- Preferred (PR):  $S$  is an extension iff it is a  $\subseteq$ -maximal complete extension,
- Stable (ST):  $S$  is an extension iff it is a conflict-free s.t.  $\forall b \in \mathcal{A} \setminus S, \exists a \in S$  s.t.  $(a, b) \in \mathcal{R}$ ,
- Grounded (GR):  $S$  is an extension iff it is the unique  $\subseteq$ -minimal complete extension.
- Ideal (ID):  $S$  is the unique ideal extension iff it is the  $\subseteq$ -maximal admissible set included in all preferred extensions
- Semi-Stable (SST):  $S$  is an extension iff  $S$  is a complete extension s.t.  $S^+$  is  $\subseteq$ -maximal.

- Stage (STG):  $S$  is an extension iff  $S$  is a conflict-free set s.t.  $S^+$  is  $\subseteq$ -maximal.

We consider both decision problems from the Approximate Track of ICCMA 2023:<sup>1</sup>

- $DC\text{-}\sigma$ : Given an AF  $\mathcal{F} = \langle \mathcal{A}, \mathcal{R} \rangle$  and  $a \in \mathcal{A}$ , is  $a$  a member of some  $\sigma$ -extension of  $\mathcal{F}$ ?
- $DS\text{-}\sigma$ : Given an AF  $\mathcal{F} = \langle \mathcal{A}, \mathcal{R} \rangle$  and  $a \in \mathcal{A}$ , is  $a$  a member of each  $\sigma$ -extension of  $\mathcal{F}$ ?

## II. APPROXIMATION APPROACH

### A. Related Work: Harper++

For all considered semantics  $\sigma$  (except STG), if an argument  $a$  belongs to the GR extension then it belongs to every  $\sigma$ -extension. If on the contrary  $a$  is attacked by an argument in the GR extension, then  $a$  does not belong to any  $\sigma$ -extension. Moreover, computing the GR extension can be done in polynomial time. So it has been proposed to use the GR extension as an approximation of other forms of reasoning in abstract argumentation. More precisely, the solver Harper++ [2] has proposed to solve all the decision problems from ICCMA 2021 as follows. For any semantics  $\sigma$  under consideration,

- Harper++ answers YES to  $DS\text{-}\sigma$  queries iff the given argument is in  $GR(\mathcal{F})$ ,
- Harper++ answers NO to  $DC\text{-}\sigma$  queries iff the given argument is attacked by an argument in the grounded extension.

Said otherwise, it always says YES (resp. NO) if the argument belongs to (resp. is attacked by) the grounded extension. The status of arguments which are not in the grounded extension nor attacked by it is assigned to “accepted” for  $DC\text{-}\sigma$  queries, and to “rejected” for  $DS\text{-}\sigma$  queries. Our goal is to propose a new approach to deal with these arguments.

### B. ARIPOTER-Degrees

Our approach consists in comparing the in-degree (*i.e.* the number of direct attackers) and the out-degree (*i.e.* the number of attackees) of an argument to decide whether it should be

accepted or not. Let us write  $\delta^-(a)$  and  $\delta^+(a)$  the in-degree and the out-degree of  $a$  respectively.

Our approach ARIPOTER-Degrees is parameterized by  $k \in \mathbb{R}$ , and  $\forall a \in \mathcal{A}$  it answers YES iff  $a$  is in the grounded extension or  $\delta^+(a) \geq k \times \delta^-(a)$ . Otherwise, the answer is NO.

Notice that we do not handle differently DC- $\sigma$  and DS- $\sigma$  queries by nature. But an experimental evaluation of the approach provides well suited values of  $k$  for the various queries.

### III. SYSTEM

#### A. Tool Description

We have implemented our approximation approach for DC- $\sigma$  and DS- $\sigma$  queries in a Java tool which is available online.<sup>2</sup>

Argumentation frameworks are represented as double adjacency lists. For a better computation of the in-degree and out-degree of arguments, each argument in the AF is associated with the list of its attackers, and the list of arguments it attacks.

The core components of the tool are the interface `ArgumentationFramework`, which provides methods to manipulate the sets of arguments and attacks, and the abstract class `Solver`, which provides the method `solve` where concrete algorithms for various reasoning tasks must be implemented. Adding a new reasoning task can thus be made by implementing a new concrete class which inherits from `Solver`, with the dedicated algorithm.

Our tool follows the specifications of ICCMA'23, regarding the command line interface, the input file format, and the standard output.

#### B. Choice of parameters

Table I describes our choice of parameter for solving the various problems of ICCMA 2023, *i.e.* the value of  $k$ . We have conducted some experiments to choose the value with provides the best accuracy on our test set.

TABLE I  
CHOICE OF PARAMETERS

Semantics $\sigma$	DC- $\sigma$	DS- $\sigma$
CO	$ \mathcal{A} $	n/a
PR	n/a	$ \mathcal{A} $
ST	$ \mathcal{A} $	0.1
ID	$ \mathcal{A} $	n/a
SST	$ \mathcal{A} $	$ \mathcal{A} $
STG	0	$ \mathcal{A} $

DC-PR is omitted since it corresponds to DC-CO. DS-CO is omitted since it corresponds to the (polynomially solvable) DS-GR, which is not included in the competition. Finally DS-ID corresponds to DC-ID.

### REFERENCES

- [1] P. M. Dung. On the acceptability of arguments and its fundamental role in nonmonotonic reasoning, logic programming and n-person games. *Artif. Intell.*, 77(2):321–358, 1995.
- [2] M. Thimm. Harper++: Using grounded semantics for approximate reasoning in abstract argumentation. <http://argumentationcompetition.org/2021/downloads/harper++.pdf>, 2021.

<sup>2</sup>[https://github.com/jeris90/approximate\\_inout](https://github.com/jeris90/approximate_inout)